

Algorithmic Desingularization

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Introduction by Pictures: Singularities

Algorithmic
Desingularization

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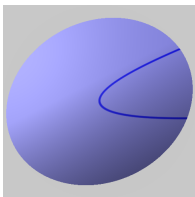
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Desingularization

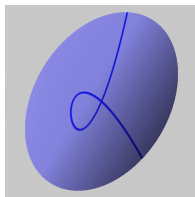
Blowing Up

Finding Centers

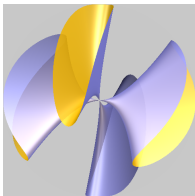
Applications



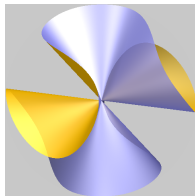
$$y^2 - x = 0$$



$$y^2 - x^2 - x^3 = 0$$



$$x^2 + y^4 - z^4 - 3x^2y^2 = 0$$



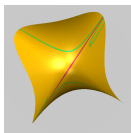
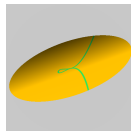
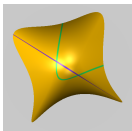
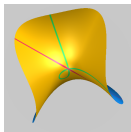
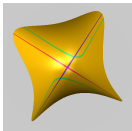
$$x^3 + y^3 - z^3 - 11xyz = 0$$

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Algorithmic
Desingularization

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Resolution of $V(y^2 + x^4 - x^5) \subset \mathbb{A}^2$



Introduction

Desingularization

Blowing Up

Finding Centers

Applications

Resolution of Singularities

Algorithmic
Desingularization

A. Frühbis-Krüger

Introduction

Desingularization

Blowing Up

Finding Centers

Applications

K field of characteristic zero
for this talk \mathbb{C}

W smooth, puredimensional scheme of dimension n
 $X \subset W$ reduced subscheme

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Algorithmic
Desingularization

A. Frühbis-Krüger

Introduction

Desingularization

Blowing Up

Finding Centers

Applications

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Simplest Formulation(non-embedded):

Find a non-singular \tilde{X} and a proper birational morphism

$$\pi : \tilde{X} \longrightarrow X$$

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such that $\text{Reg}(X) \cong \pi^{-1}(\text{Reg}(X))$.

Blowing up - in Pictures

Algorithmic
Desingularization

A. Frühbis-Krüger

Introduction

Desingularization

Blowing Up
Finding Centers

Applications

Idea: Replace a point in the plane
by a projective line

Blowing up - in Pictures

Algorithmic
Desingularization

A. Frühbis-Krüger

Introduction

Desingularization

Blowing Up
Finding Centers

Applications

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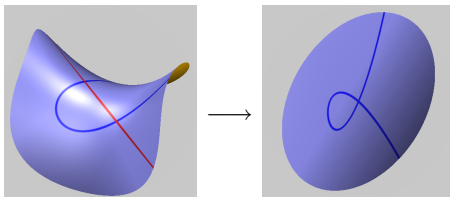
effect: more room for curves to become smooth

Blowing up - in Pictures

Idea: Replace a point in the plane
by a projective line

effect: more room for curves to become smooth

in pictures (just 1 chart):



Embedded Desingularization

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Desingularization

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Technical Formulation:

Find a finite sequence of blowing-ups

$$W_r \xrightarrow{\pi_r} \cdots \xrightarrow{\pi_2} W_1 \xrightarrow{\pi_1} W_0 = W$$

at smooth centers $C_i \subset W_i$ such that

Introduction

Desingularization

Blowing Up

Finding Centers

Applications

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5. $(W_r, X_r) \longrightarrow (W, X)$ is equivariant under group action

History of the problem

Algorithmic
Desingularization

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Introduction

Desingularization

Blowing Up
Finding Centers

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- ▶ Curves: L.Kronecker, M.Noether, A.Brill, ... (1890s)

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approaches of more algebraic flavor:

- ▶ surfaces/3-folds: O.Zariski (1930s/40s)
- ▶ general case: **H.Hironaka (1964)** for $\text{char}K = 0$

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recent developments:

- ▶ algorithmic proofs: Bierstone+Milman; Villamayor; Encinas+Hauser (since 1990s)
- ▶ implementations: Bodnar+Schicho; FK+Pfister

Main Algorithmic Tasks

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Introduction

Desingularization

Blowing Up

Finding Centers

Applications

For each blowing up step:

Main Algorithmic Tasks

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A Blowing up of W_i along a given center C_i

- ▶ Groebner basis computation in at least $(n + \text{codim}(C_i) + 1)$ variables
- ▶ iterated ideal quotients (=again Groebner bases)
- ▶ algorithmically straight forward, not overly expensive

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B Finding suitable centers C_i

key difficulty: not all permissible centers improve the situation

Computing a blowing up

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Desingularization

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Introduction

Desingularization

Blowing Up
Finding Centers

Applications

X affine variety, its ideal $I_X \subset K[x_1, \dots, x_n]$

C smooth subvariety of X

(wlog $I_C = \langle f_1, \dots, f_k \rangle$)

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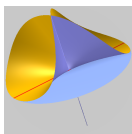
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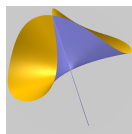
can be computed as preimage of I_X under

$$\begin{aligned} \Phi : K[x_1, \dots, x_n, y_1, \dots, y_k] &\longrightarrow K[x_1, \dots, x_n, t] \\ x_i &\longmapsto x_i \\ y_j &\longmapsto t \cdot f_j \end{aligned}$$

Some Good and Bad Choices of Centers I



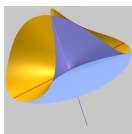
$C=0$
 \longrightarrow



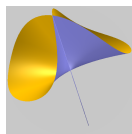
$$U_x: X_{strict} = V(z^2 + xy^2), \\ E = V(x)$$

$$X = V(z^2 + xy^2)$$

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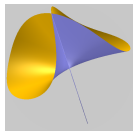


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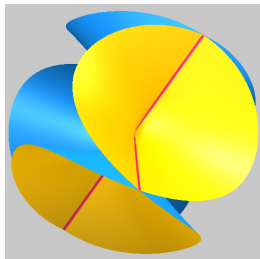
$$X = V(z^2 + xy^2)$$

smooth

$$\xrightarrow{C=line}$$

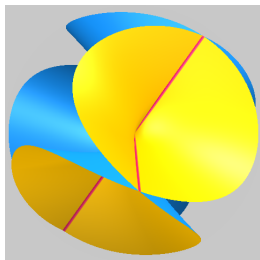


Some Good and Bad Choices of Centers II



$$X = V(z^2 - x^2y^2) \subset \mathbb{A}^3$$

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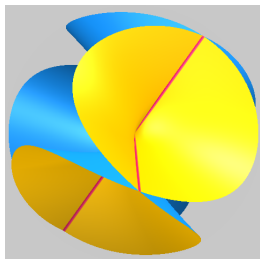


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Choices of Center:

- ▶ $\text{Sing}(X)$ singular \implies impossible

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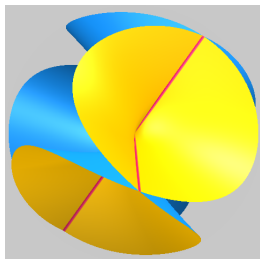


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- ▶ $\mathbf{0}$ is only possible choice

General Philosophy for Finding the Center

Algorithmic
Desingularization

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Introduction

Desingularization

Blowing Up

Finding Centers

Applications

'worst' points are points of maximal value of a governing
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General Philosophy for Finding the Center

Algorithmic
Desingularization

A. Frühbis-Krüger

Introduction

Desingularization

Blowing Up

Finding Centers

Applications

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Algorithmic
Desingularization

A. Frühbis-Krüger

Introduction

Desingularization

Blowing Up

Finding Centers

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Algorithmic
Desingularization

A. Frühbis-Krüger

Introduction

Desingularization

Blowing Up

Finding Centers

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- ▶ maximal value does not increase under blowing ups
- ▶ decrease of maximal value measures progress of desingularization

General Structure of Invariant

$$(inv_n; inv_{n-1}; \dots; inv_2)$$

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- ▶ depending on algorithmic approach inv_n can also have the Hilbert-Samuel function at w as first entry

Variants in SINGULAR

Variants of embedded desingularization:

- ▶ variant of Villamayor's Algorithm (available):
 - + all dimensions
 - + no special conditions on ideal
 - large amount of data
 - a number of unproductive blowing ups

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 - even more data due to further splitting up of charts
- ▶ Blanco's variant for binomial ideals (implemented, not yet in distribution):
 - + computation of center by combinatorial process
 - + faster, total amount of data smaller
- ▶ Jung's algorithm for surfaces (implementation FK-Renner):
 - only for surfaces
 - + faster, fewer charts

Problems/Tasks

result of resolution process represented in charts

⇒ need to extract desired information

Algorithmic
Desingularization

A. Frühbis-Krüger

Introduction

Desingularization

Blowing Up

Finding Centers

Applications

Problems/Tasks

Algorithmic
Desingularization

A. Frühbis-Krüger

Introduction

Desingularization

Blowing Up

Finding Centers

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- ▶ Identification of exceptional divisors in different charts

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- ▶ Identification of subvarieties in different charts
- ▶ Identification of exceptional divisors in different charts
- ▶ Separation of \mathbb{C} -irreducible components of exceptional divisors

Applications of desingularization

Given $\pi : W \longrightarrow \mathbb{C}^n$ embedded resolution of $V = f^{-1}(0)$,
 E_i irreducible components of $\pi^{-1}(f^{-1}(0))$
 N_i multiplicity of E_i in divisor of $f \circ \pi$
 $\nu_i - 1$ multiplicity of E_i in divisor of $\pi^*(dx_1 \wedge \cdots \wedge dx_n)$

Currently available:

- ▶ intersection form of exceptional curves on desingularized surface
- ▶ dual graph of resolution (surface case)
- ▶ discrepancies $a_i = \nu_i - N_i$
- ▶ topological zeta function (global and local)

$$Z_{top,f}(s) = \sum_I \chi(E_I^\circ) \prod_{i \in I} \frac{1}{N_i s + \nu_i}$$